

S-72.245 Transmission Methods in Telecommunication Systems (4 cr)

Sampling and Pulse Coded Modulation

Sampling and Pulse Coded Modulation

- Pulse amplitude modulation
- Sampling
 - Ideal sampling by impulses
 - practical chopper sampler
- Line coding
- Quantization
 - Uniform
 - Non-uniform
 - μ -law - compression
 - quantization noise
- PCM and channel noise
- PCM multiplexing

TDM: Time Division Multiplexing
FDM: Frequency Division Multiplexing
PAM: Pulse Amplitude Modulation
PCM: Pulse Coded Modulation

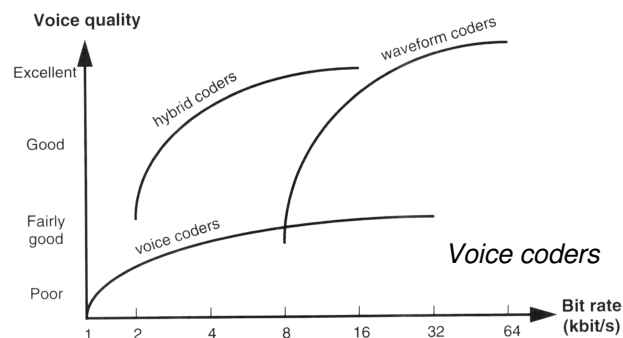
Short history of pulse coded modulation

- A problem of PSTN analog techniques (eg SSB-FDM) was that transmitting multiple channels was difficult due to non-linearities resulting channel **cross-talk**
- **1937** Reeves and Delorane ITT labs. tested **TDM**-techniques by using electron-tubes
- **1948 PCM** tested in Bell Labs: Using this method it is possible to represent a 4 kHz analog telephone signal as a **64 kbit/s digital bit stream**
- TDM was taken into use in **1962** with a **24 channel PCM link**
- The first **30-channel** PCM system installed **in Finland 1969**
- Nowadays all exchanges in Finland use ISDN & PCM based cables, microwave or optical links

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PCM coding is a form of waveform coding



- **Waveform coders** reply signal by quantized (discrete) values, - precise waveform replay but requires a lot of bandwidth
- **Parameterized coders** count on **system model** that reproduces the signal. Only **model parameters** are transmitted and updated. Very low rate can be obtained but this is paid by quality degradation
- **Hybrid coders** (as Δ - modulation) are a compromise solution

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Some important ITU-T speech/video coding standards

Standard no. Name	Description	Current status
G.711	Pulse code modulation (PCM) of voice frequencies (64 kbit/s)	Adopted 1984
G.722, G.725	7 kHz audio-coding within 64 kbit/s	Adopted 1988
G.726	16/24/32/46 kbit/s adaptive differential pulse code modulation (ADPCM)	Adopted 1990
G.728	16 kbit/s speech coding with excited linear prediction	Adopted 1992
G.729	8 kbit/s speech coding	Adopted 1996
H.221	Frame structure for a 64 to 1920 kbit/s channel in audiovisual teleservices	Adopted 1990
H.230	Control and indication signals for audiovisual systems	Adopted 1990
H.231, H243	Multipoint videoconferencing	Adopted 1993
H.233	Encryption / Privacy systems	Adopted 1993
H.261	Video codec for audiovisual teleservices at p x 64 kbit/s	Adopted 1993
H.263	Video coding for low bit rate communication	Adopted 1996
MPEG1	Stored motion video stored at <2 Mbit/s	Adopted 1993
MPEG2	Stored/live motion video at 5–60 Mbit/s	Adopted 1994
MPEG4	Low bit rate (<64 kbit/s) coding of motion video	Adopted 1991
JPEG	Still-frame graphics for multimedia	Adopted 1991

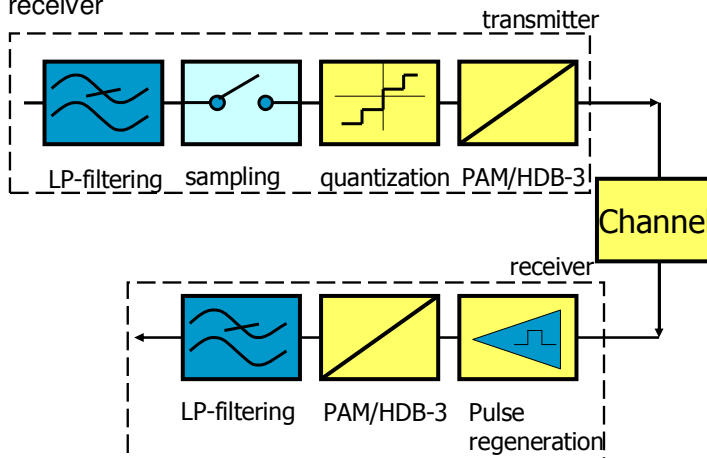
MPEG = Motion picture expert group. JPEG = Joint photograph expert group

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Pulse Coded Modulation (PCM)

- PCM is a method by which an analog message can be transformed into numerical format and then decoded at the receiver



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PCM principles: Ideal sampling

- The rectangular pulse train

$$s(t) = \sum_{k=-\infty}^{\infty} \Pi\left(\frac{t - kT_s}{\tau}\right)$$

- The ideal sampling pulse train

$$s_s(t) = \sum_{k=-\infty}^{\infty} \lim_{\tau \rightarrow 0} \left[\frac{1}{\tau} s(t - kT_s) \right] = \sum_{k=-\infty}^{\infty} \delta(t - kT_s)$$

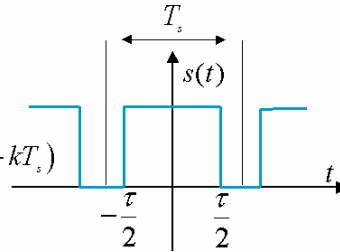
- The ideally sampled signal is a pulse train of weighted impulses

$$\begin{aligned} x_s(t) &= x(t)s_s(t) = x(t) \sum_{k=-\infty}^{\infty} \delta(t - kT_s) \\ &= \sum_{k=-\infty}^{\infty} x(kT_s) \delta(t - kT_s) \end{aligned}$$

- Translation Fourier tables: $\mathcal{F}\left[\sum_{k=-\infty}^{\infty} \delta(t - kT_s)\right] = f_s \sum_{n=-\infty}^{\infty} \delta(f - nf_s)$

the ideally sampled signal is then $X_s(f) = X(f) \otimes s_s(f)$

$$X_s(f) = f_s \sum_{n=-\infty}^{\infty} X(f - nf_s) = X(f) \otimes f_s \sum_{n=-\infty}^{\infty} \delta(f - nf_s)$$



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Ideal sampling: reconstruction

- Reconstruction is obtained by lowpass (LP) filtering. Assume the ideal LP filter with

$$H(f) = K \Pi\left(\frac{f}{2B}\right) \exp(-j\omega t_d)$$

- Due to the translations

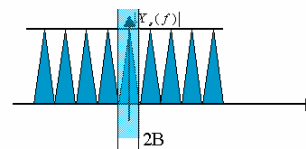
$$\text{sinc } 2Wt \leftrightarrow \frac{1}{2W} \Pi\left(\frac{f}{2W}\right) \quad v(t - t_d) \leftrightarrow V(f) \exp(-j\omega t_d)$$

the respective impulse response of the ideal LP is therefore

$$h(t) = 2BK \text{sinc } 2B(t - t_d)$$

- In ideal sampling reconstruction weighted impulse train (representing the sampled signal) is applied to this filter and the output is

$$\begin{aligned} y(t) &= h(t) \otimes x_s(t) \\ &= 2BK \text{sinc } 2B(t - t_d) \otimes \sum_{k=-\infty}^{\infty} x(kT_s) \delta(t - kT_s) \\ &= 2BK \sum_{k=-\infty}^{\infty} x(kT_s) \text{sinc } 2B(t - t_d - kT_s) \end{aligned}$$



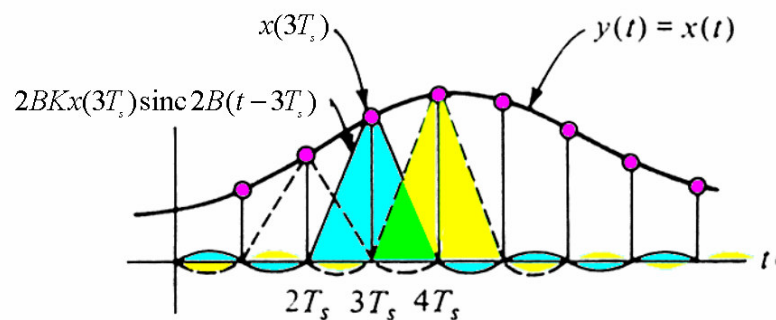
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Reconstructed time-domain signal consists of interpolated sinc-functions

- At the sample instances all but one sinc functions are zero
- Therefore all *band limited* signals can be expressed by using the sinc-series:

$$2BK \sum_{k=-\infty}^{\infty} x(kT_s) \text{sinc } 2B(t - t_d - kT_s)$$



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Aliasing and sampling theorem

- Nyquist sampling theorem:

If a signal contains no frequency components for $|f| \geq W$ it is **completely** described by instantaneous uniformly spaced time samples having period $T_s \leq 1/2W$. The signal can hence be reconstructed from its samples by an ideal LPF of bandwidth B such that $W \leq B \leq f_s - W$

- Note: If the signal contains higher frequencies than twice the sampling frequency they **will also be present** at the sampled signal! An application of this is the sampling oscilloscope (next slide)

- Also, it follows from the sampling theorem that

Two pieces of independent information / second (independent samples) can be transmitted in 1 Hz wide channel

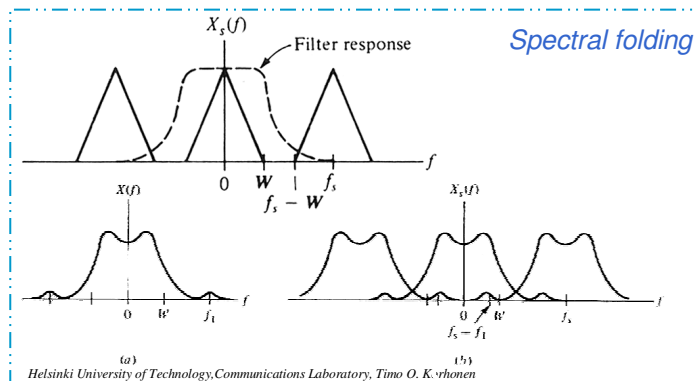
because signal having bandwidth B can be constructed from rate $2B$ independent samples

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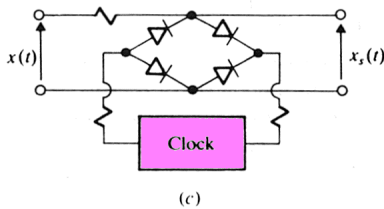
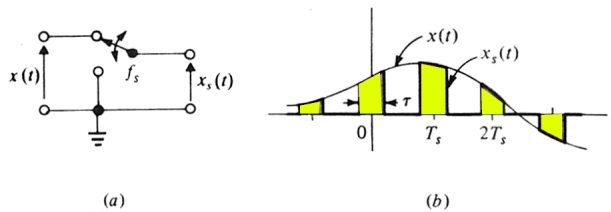
Unperfect reconstruction

1. Sampling wave pulses have finite duration and risetimes -> linear distortion
2. Reconstruction filters are not ideal lowpass filters -> spectral folding
3. Sampled messages are time limited and therefore their spectra is not frequency limited -> spectral folding
4. Samples digitized by finite length words -> quantization noise

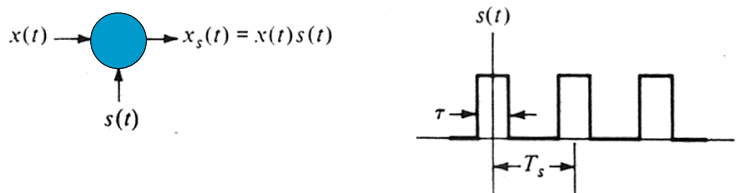


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Chopper sampling



The chopper sampler waveforms

- Sampling wave consists of a periodic pulse train whose duration is τ and period is T_0

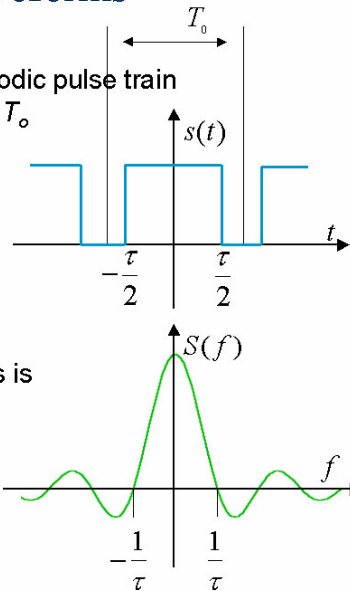
$$c_n = \frac{1}{T_0} \int_{-\tau/2}^{\tau/2} v(t) \exp(-j2\pi n t / T_0) dt$$

$$c_n = \frac{A\tau}{T_0} \text{sinc}(nf_0\tau)$$

$$s(t) = \sum_{n=-\infty}^{\infty} c_n \exp(j2\pi n f_0 t)$$

The Fourier series for real signals is

$$s(t) = c_0 + \sum_{n=1}^{\infty} 2c_n \cos n\omega_s t$$



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Chopper sampler: sampled spectra

- Consider the sampled signal from the chopper sampler by term-by-term multiplication

$$x_s(t) = x(t)s(t)$$

$$= x(t) \left[c_0 + \sum_{n=1}^{\infty} 2c_n \cos n\omega_s t \right]$$

$$x_s(t) = x(t)s(t)$$

$$= x(t)c_0 + 2c_1 x(t) \cos \omega_s t + 2c_2 x(t) \cos 2\omega_s t \dots$$

Remember the modulation theorem:

$$v(t) \cos(\omega_c t + \phi) \leftrightarrow \frac{1}{2} [V(f - f_c) \exp(j\phi) + V(f + f_c) \exp(-j\phi)]$$

- Therefore the sampled signal is in frequency domain

$$X_s(f) = c_0 X(f) + c_1 [X(f - f_s) + X(f + f_s)] \\ + c_2 [X(f - 2f_s) + X(f + 2f_s)] \dots$$

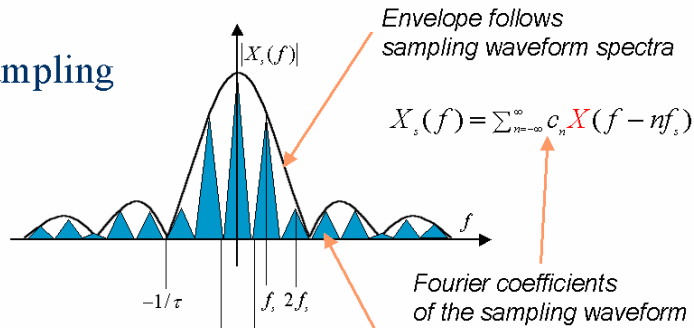
$$s(t) = c_0 + \sum_{n=1}^{\infty} 2c_n \cos n\omega_s t$$

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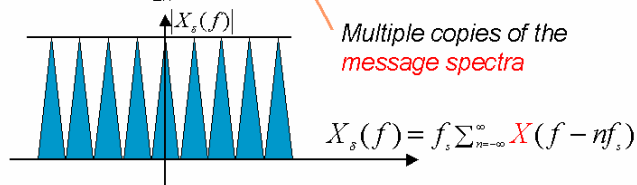
Spectra of ideal sampling and chopper sampling compared

Chopper sampling



Ideal sampling

Message's **baseband** spectral width equals W



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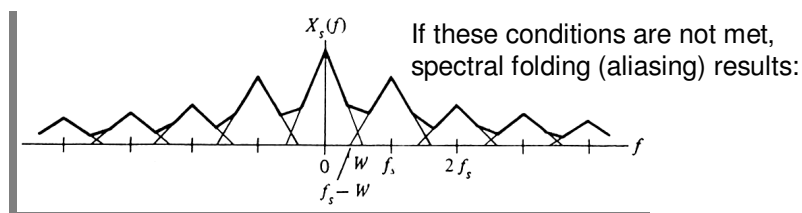
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Observations on chopper sampling

- Resulting spectra
 - has the envelope of the sampling waveform
 - has the sampled signal repeated at the integer multiples of the sampling frequency
- Therefore the sampled signal can be reconstructed by filtering provided that

$$X(f) = 0, |f| > W \quad \text{Sampled signal is band limited}$$

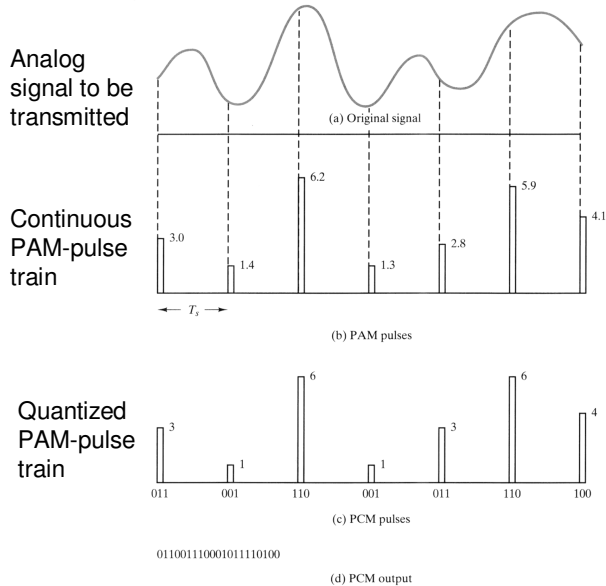
$$f_s \geq 2W \quad \text{Sampling rate is high enough}$$



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Quantization

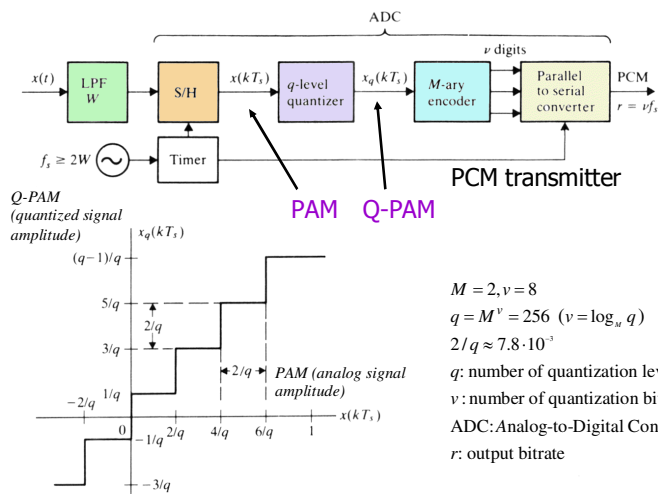


- Original signal has continuous amplitudes in its dynamic range
- PAM - signal is a discrete constant period, pulse train having continuous amplitude values
- Quantized PAM signal has only the values that can be quantized by the words available (here by 3 bit words)

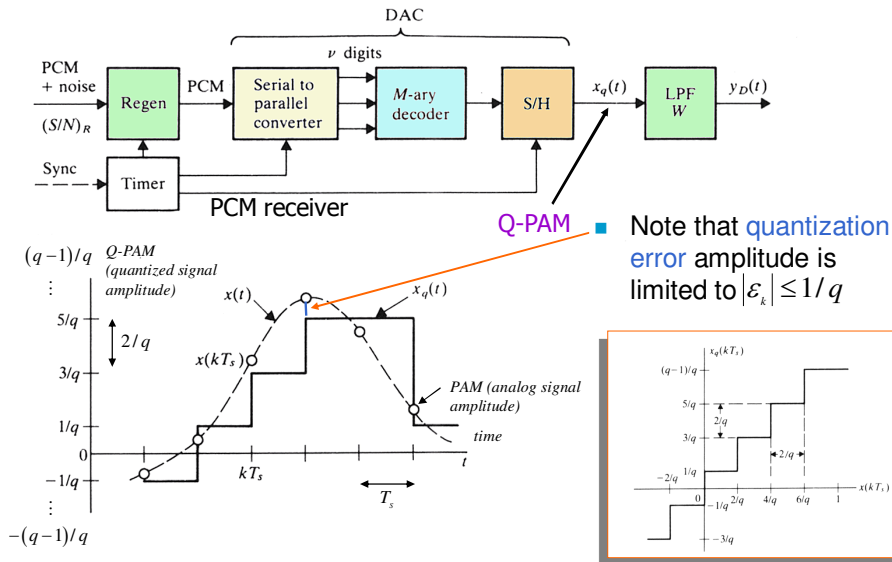
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Uniform quantization: transmitter

- Transforming the continuous samples into discrete level samples is called quantization
- In uniform quantization quantization step size is constant



Reconstruction from the quantized signal



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Quantization noise: uniform quantization

- Model the quantized signal by assuming ideal PAM sampling using the quantization error ϵ_k :

$$y(t) = \sum_k [x(kT_s) + \epsilon_k] \delta(t - kT_s)$$

- Quantization error is the difference of the PAM and the Q-PAM signals

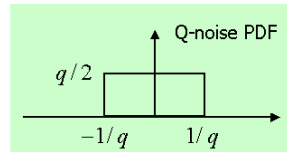
$$\epsilon_k = x_q(kT_s) - x(kT_s)$$

- The final output is obtained by using the ideal LPF:

$$y_D(t) = x(t) + \sum_k \epsilon_k \text{sinc}(f_s t - k)$$

- Assuming equal probable signal at all amplitude levels yields for quantization noise average power

$$\overline{\sigma_\epsilon^2} = \frac{q}{2} \int_{-1/q}^{1/q} \epsilon^2 d\epsilon = \frac{1}{3q^2}$$



note: $\frac{2}{q} \frac{q}{2} = 1$

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PCM versus analog modulation

- PCM suppresses wideband noise
- A PCM system well above threshold

$$S_D / N_D = 3q^2 S_x \approx 3M^{2b} S_x, q = M^n$$

$$\text{transmission BW: } B_T = bW$$

- Observations:
 - operate PCM near threshold!
- Note that for radio 75 dB dynamic range (12 bits without compression) transmission bandwidth B_T is

$$W = 15 \text{ kHz}, n = 12 \text{ bit}$$

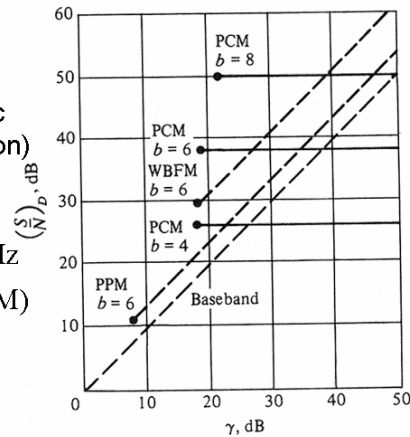
$$B_T \approx nf_s / 2 = 12 \times 35 / 2 = 210 \text{ kHz}$$

(This is about the same as in FM)

- Why to use PCM then??
 - regenerative digital repeaters
 - digital multiplexing
 - etc!

$$\left(\frac{S_D}{N_D} \right)_{PM} = 3D^2 S_x \gamma$$

$$\left(\frac{S_D}{N_D} \right)_{PCM} = \frac{3q^2}{1 + 4q^2 P_e} S_x$$



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Line coding

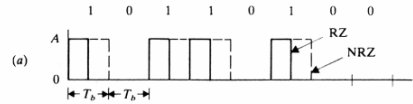
- Line codes are used to **enable baseband transmission** in
 - Fiber optic systems
 - Cable transmission
 - Data access and storage
- In PCM-links line coding is used to alleviate clock **synchronization** at the receiver (F-transform of the pulse train should contain spikes that the receiver clock can be synchronized)
- Line codes should
 - be immune to **long strings of zeros** that can lead to missing receiver clock synchronization
 - contain zero long term averaged **DC-component**
 - have minimum **bandwidth**
- Line codes can also be used for **error detection**

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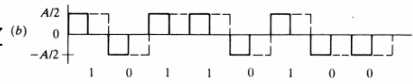
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Line coding (cont.)

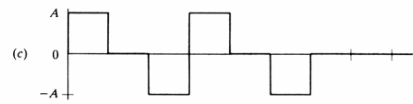
Unipolar [0,A] RZ and NRZ



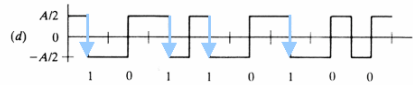
Polar [-A/2,A/2] RZ and NRZ



Bipolar [-A/2,0,A/2] AMI



Split-Phase Manchester

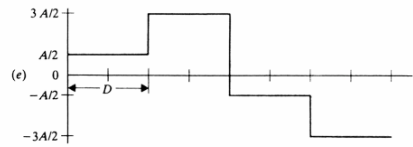


Code rate
reduced
by n

Split-Polar quaternary NRZ

$$M = 2^n$$

$$r_{se} = r/n = r/\log_2 M$$



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Line coding methods ... (cont.)

Bipolar-AMI

- 0 = no line signal
- 1 = positive or negative level, alternating for successive ones

Pseudoternary

- 0 = positive or negative level, alternating for successive zeros
- 1 = no line signal

Manchester

- 0 = transition from high to low in middle of interval
- 1 = transition from low to high in middle of interval

Differential Manchester

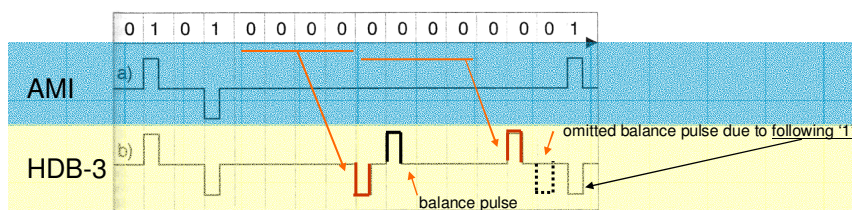
- Always a transition in middle of interval
- 0 = transition at beginning of interval
- 1 = no transition at beginning of interval

B8ZS

- Same as bipolar AMI, except that any string of eight zeros is replaced by a string with two code violations

HDB3

- Same as bipolar AMI, except that any string of four zeros is replaced by a string with one code violation



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Repeaters

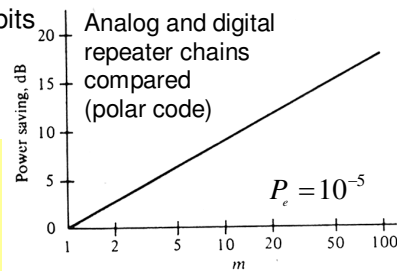
- At the transmission path regenerative repeaters are often used
- At the receiver signal is transformed back to analog form by lowpass filtering removing harmonics produced by sampling
- Repeaters are categorized as:
 - **analog repeater**: gain equal to the line attenuation between repeaters
 - **digital repeater**: regenerates bits by decoding and encoding

Error rates for polar, baseband, m -stage repeater chains:

$$S_d / N_d \approx \frac{(S/N)_1}{m} \text{ analog repeaters*}$$

$$\Rightarrow P_e = Q\left[\sqrt{(S/N)_1 / m}\right]$$

$$P_e \approx mQ\left[\sqrt{(S/N)_1}\right] \text{ digital repeaters}$$



* Formula follows from the reverse inspection than coherent averaging

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Digital Repeater Chain Error Rate

- Assume each repeater has the bit error rate

$$\alpha = Q\left[\sqrt{(S/N)_1}\right]$$

- There are m repeaters that introduce errors that are produced for *odd* number of error digits* and the *probability of i errors* in the chain is

$$P(i, m) = \binom{m}{i} \alpha^i (1-\alpha)^{m-i}, i \text{ odd}$$

$$P_e = \binom{m}{1} \alpha (1-\alpha)^{m-1} + \binom{m}{3} \alpha^3 (1-\alpha)^{m-3} \dots$$

$$\approx m\alpha \quad (\alpha \ll 1, m \text{ moderate})$$

- And therefore

$$P_e \approx mQ\left[\sqrt{(S/N)_1}\right] \left(\binom{m}{1} \alpha (1-\alpha)^{m-1} \approx \frac{m!}{(m-1)!} \alpha = m\alpha \right)$$

$$\left(\binom{m}{3} \alpha^3 (1-\alpha)^{m-3} \approx \frac{m!}{3!(m-3)!} \alpha^3 \ll m\alpha \right)$$

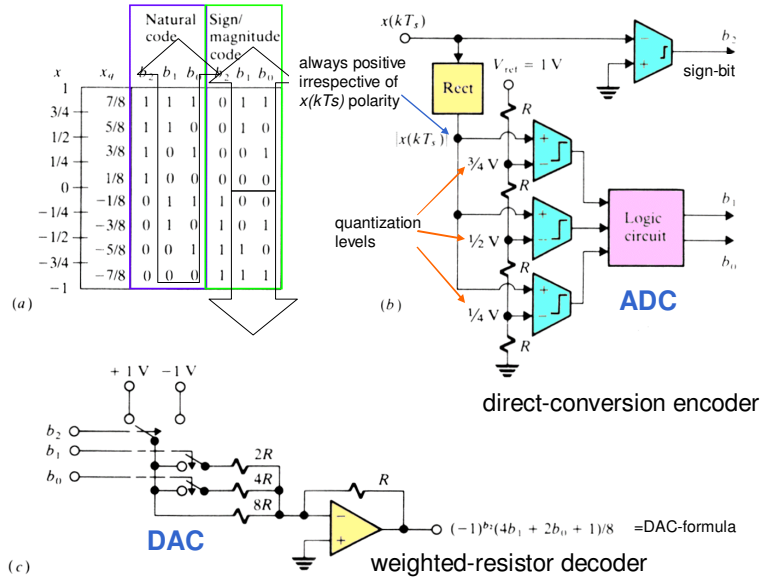
$$\approx m\alpha \quad (\alpha \ll 1, m \text{ moderate})$$

*Even number of errors cancel out each other!

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PCM encoding and decoding circuits ($n=3$)

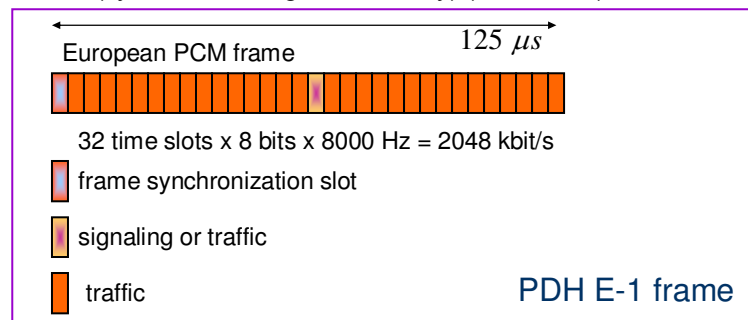


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PCM Systems and Digital Time Division Multiplexing (TDM)

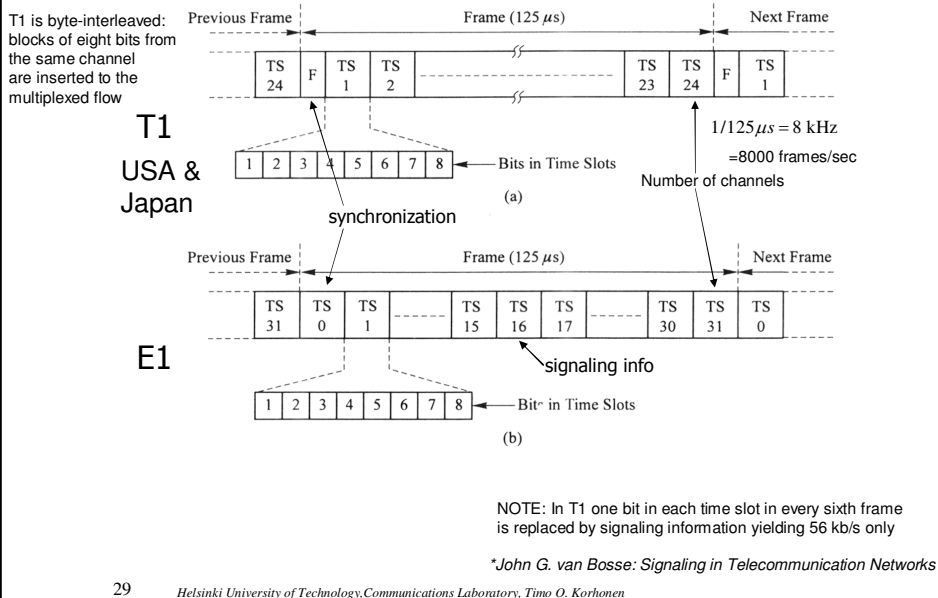
- In digital multiplexing several messages are transmitted via same physical channel. For multiplexing 64 kbit/s channels in digital exchanges following three methods are available:
 - **PDH** (plesiochronous digital hierarchy) (the dominant method today, E1 & T1) ('50-'60, G.702)
 - **SONET** (synchronous optical network) ('85)
 - **SDH** (synchronous digital hierarchy) (CCITT '88)



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E1 and T1 First Order Frames Compared*

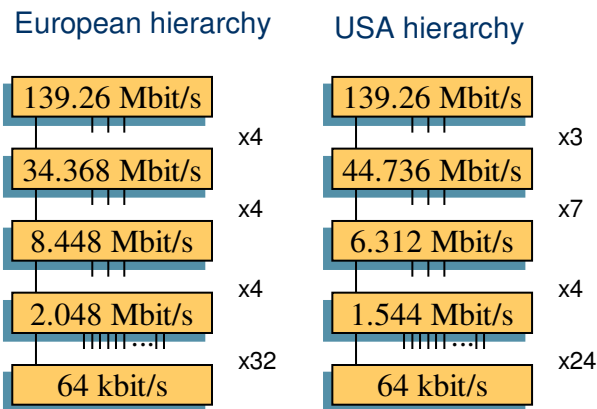


T1 and E1 Summarized

- In PSTN two PCM systems dominate:
 - T1, developed by Bell Laboratories, used in USA & Japan
 - E1, developed by CEPT* used in most of the other countries
- In both data streams divided in frames of 8000 frames/sec
- In T1
 - 24 time-slots and a framing (F) bit serves 24 channels
 - Frame length: $1 + 8 \times 24 = 193$ bits
 - Rate 193×8000 bits/second = **1544 kb/s**
- In E1
 - frame has 32 time-slots, TS 0 holds a synchronization pattern and TS 16 holds signaling information
 - An E1 frame has $32 \times 8 = 256$ bits and its rate us $8000 \times 256 =$ **2048 kb/s**

*European Conference of Postal and Telecommunications Administration

PCM Hierarchy in PDH



If one wishes to disassemble a tributary from the main flow the main flow must be demultiplexed step by step to the desired main flow level in PDH.

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PCM-method summarized

- Analog speech signal is applied into a LP-filter **restricting its bandwidth** into 3.4 kHz
- Sampling circuit forms a **PAM pulse train** having rate of 8 kHz
- Samples are **quantized** into 256 levels that requires a 8 bit-word for each sample ($2^8=256$).
- Thus a telephone signal requires $8 \times 8 \text{ kHz} = 64 \text{ kHz bandwidth}$
- The samples are **line coded** by using the HDB-3 scheme to enable synchronization and channel adaptation
- Usually one transmits several channels simultaneously following a **digital hierarchy** (as SDH or PDH)
- Transmission link can be an **optical fiber, radio link or an electrical cable**
- At the receiver the PAM signal is reconstructed where after it is **lowpass filtered** to yield the original-kind, analog signal

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