GSM

Global System for Mobile communication

GPRS

General Packet Radio Service

Examples of digital wireless systems (all originally specified by ETSI)

GSM (Global System for Mobile communication) is a *cellular mobile* system

- cellular concept
- high mobility (international roaming)

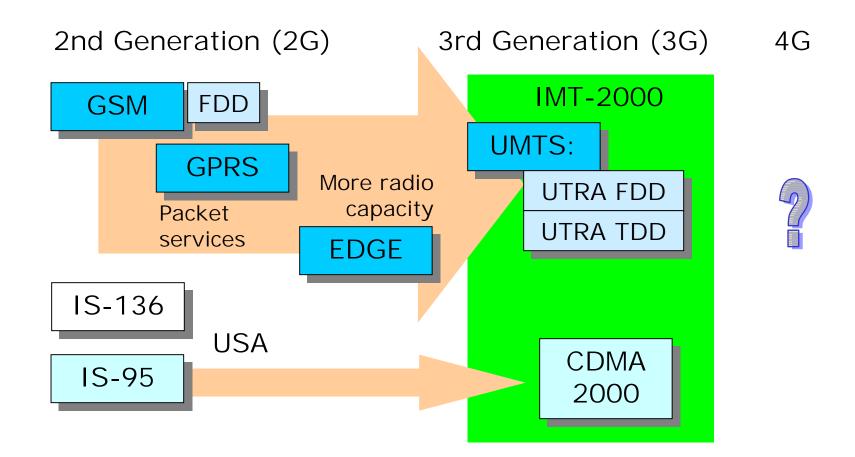
TETRA (TErrestrial Trunked RAdio) is an example of a *Professional/Privat Mobile Radio* (PMR) system

- limited access (mainly for professional usage)
- limited mobility (but other advanced features)

DECT (Digital Enhanced Cordless Telecommunications) is a *cordless* system

low mobility (only within "isolated islands")

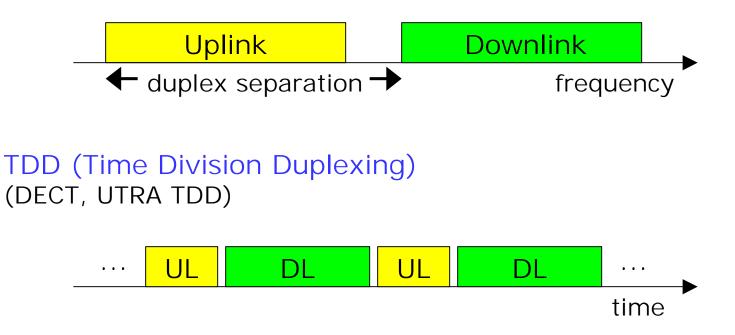
Digital PLMN systems (status 2004) (PLMN = Public Land Mobile Network)



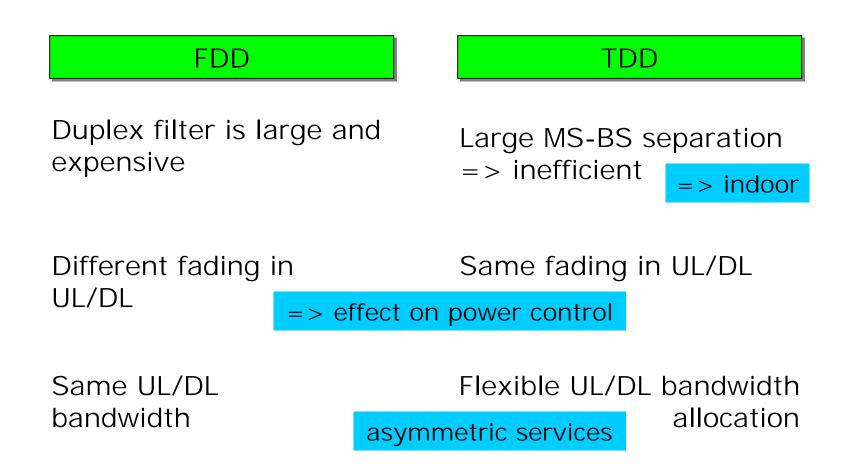
Duplexing

(separation of uplink/downlink transmission directions)

FDD (Frequency Division Duplexing) (GSM/GPRS, TETRA, UTRA FDD)

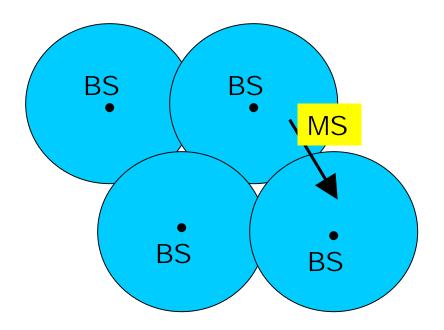


FDD vs. TDD



GSM => cellular concept

The GSM network contains a large number of cells with a base station (BS) at the center of each cell to which mobile stations (MS) are connected during a call.

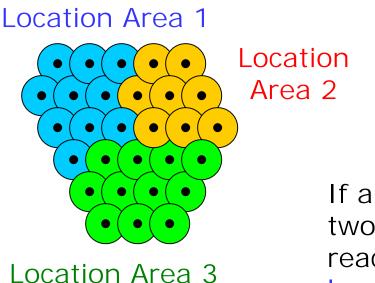


If a connected MS (MS in call phase) moves between two cells, the call is not dropped.

Instead, the network performs a handover (US: hand-off).

GSM => mobility concept

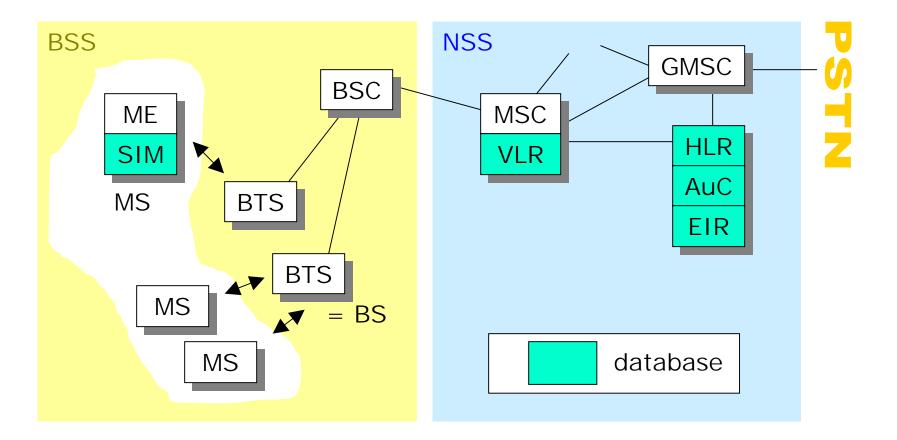
The GSM network is divided into location areas (LA), each containing a certain number of cells.



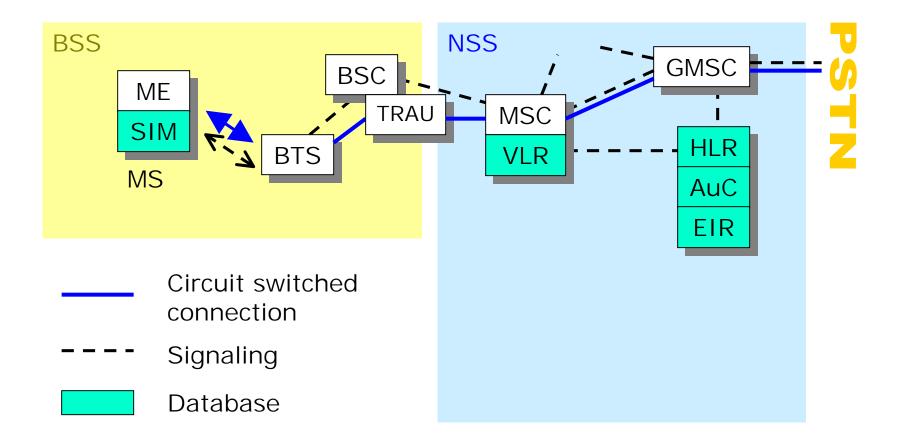
As long as an idle MS (idle = switched on) moves within a location area, it can be reached through paging.

If an idle MS moves between two location areas, it cannot be reached before it performs a location update.

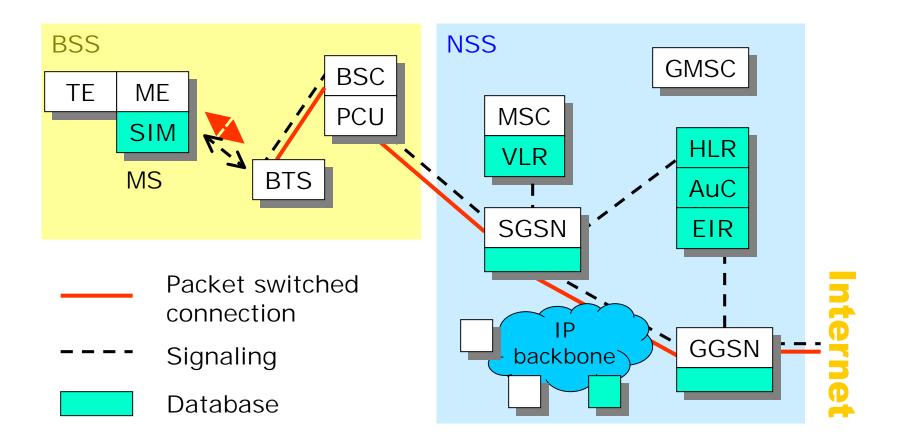
Original GSM system architecture



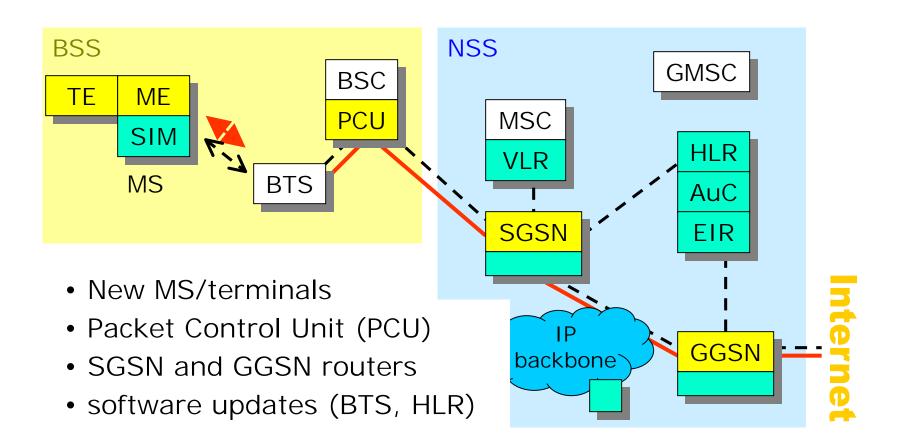
GSM: circuit switched connections



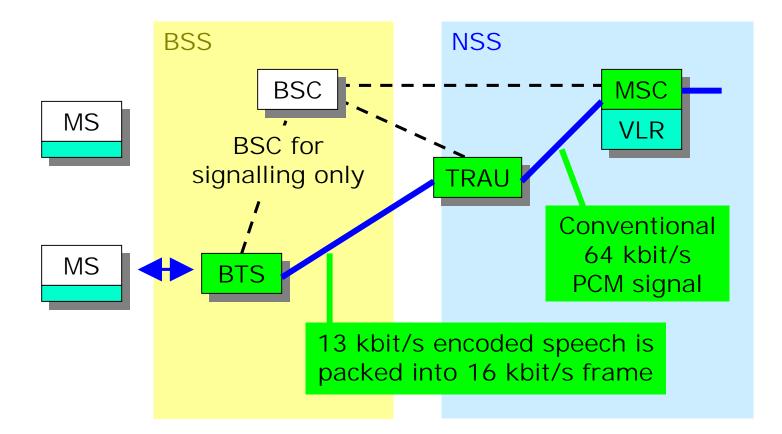
GPRS: packet switched connections



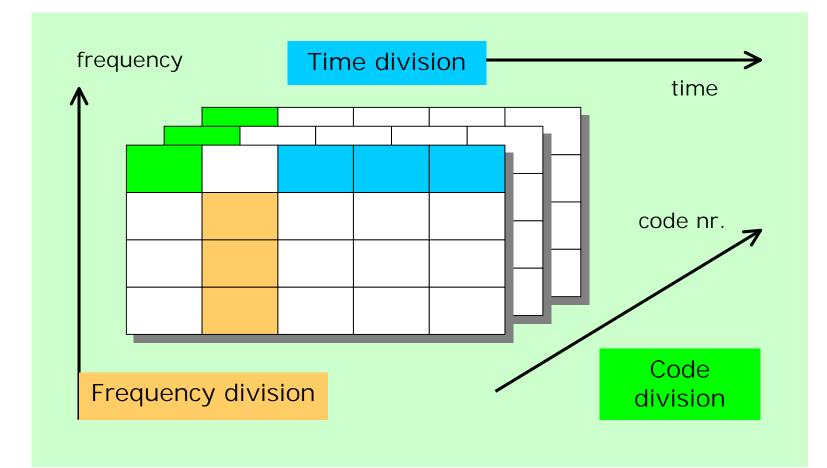
Upgrading from GSM to GSM/GPRS



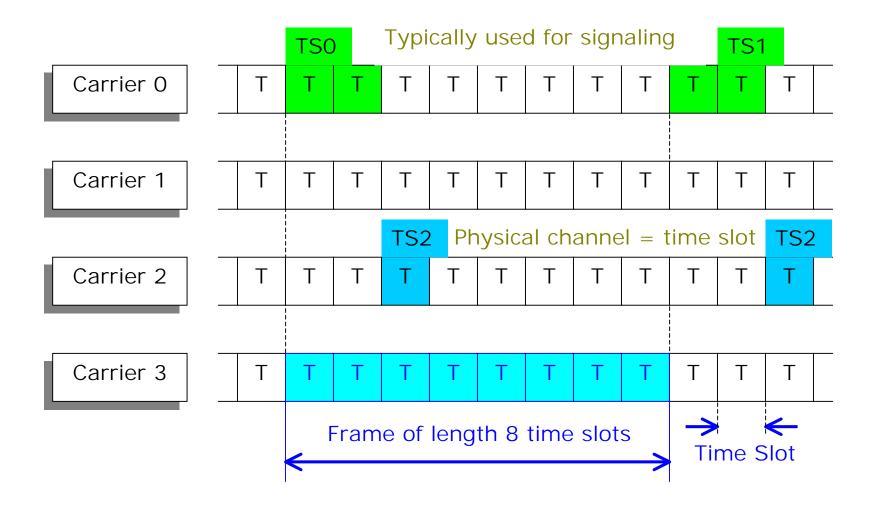
Purpose of TRAU (TRAU = Transcoding and Rate Adaptation Unit)



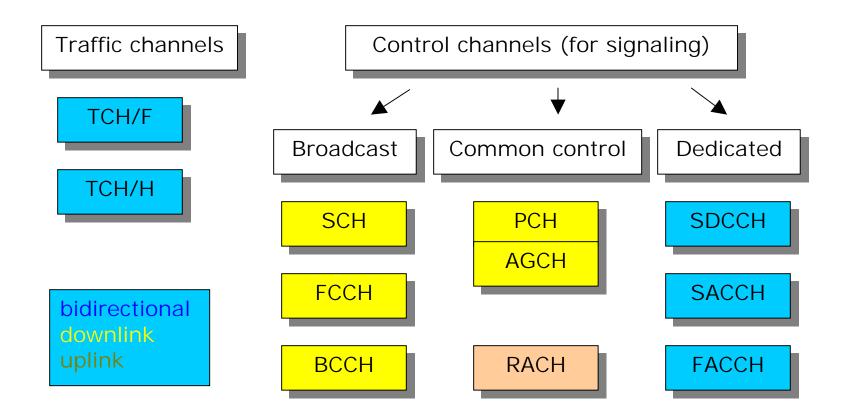
Radio interface - multiple access techniques



Physical channel = repetitive time slot

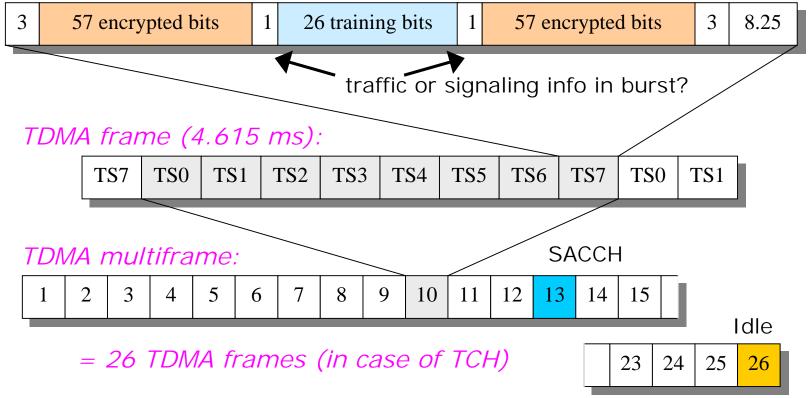


GSM logical channels



GSM burst structure

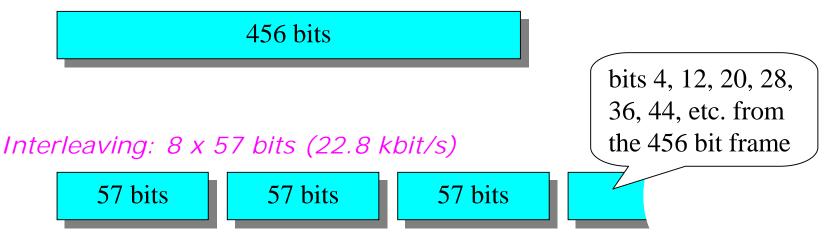
GSM normal burst: 156.25 bits (0.577 ms)



GSM speech encoding

Voice coding: 260 bits in 20 ms blocks (13 kbit/s) MS - TRAU

Channel coding: 456 coded bits (22.8 kbit/s) MS - BTS

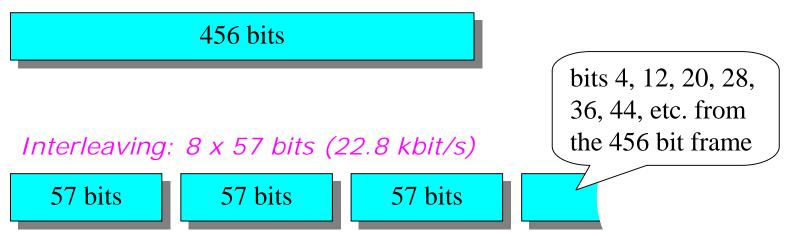


GSM signaling message encoding

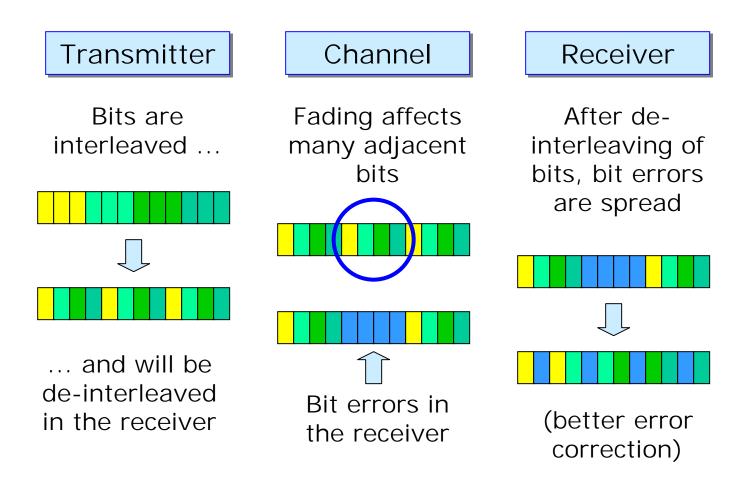
Signaling message is segmented into blocks of 184 bits:

184 bits

Each block is coded into 456 bits (22.8 kbit/s)



Purpose of interleaving



Task Management in GSM/GPRS

Radio Resource Management (RM)

- 1) Random access and channel reservation Handover management
- Ciphering (encryption) over radio interface

Number refers to the remaining slides

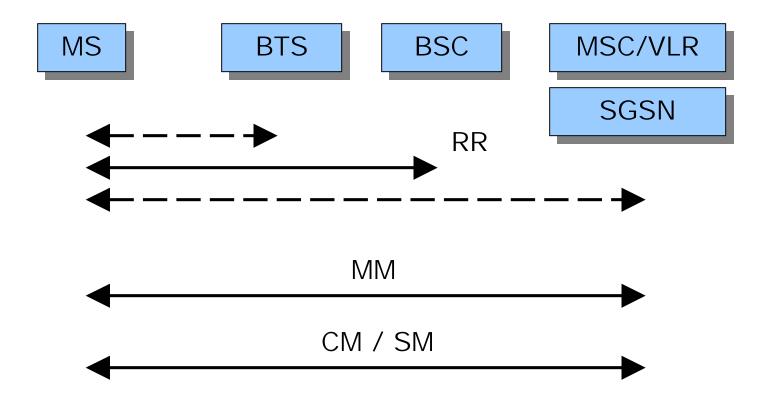
Mobility Management (MM)

4

IMSI/GPRS Attach (switch on) and Detach (switch off) Location updating (MS moves to other Location/Routing Area) Authentication

Call Control (CC) in GSMMOC, MTC5Session Management (SM) in GPRSPDP Context6

Who is involved in what?



1 Random access in GSM/GPRS (1)

Communication between MS and network is not possible before going through a procedure called *random access*.

Random access must consequently be used in

network originated activity

• paging, e.g. for a mobile terminated call in GSM

MS originated activity

- IMSI attach, IMSI detatch
- GPRS attach, GPRS detach
- location updating in GSM or GPRS
- mobile originated call in GSM
- SMS (short message service) message transfer

1 Random access in GSM/GPRS (2)

1. MS sends a short access burst over the Random Access CHannel (RACH) in uplink using Slotted Aloha (in case of collision => retransmission after random time)

2. After detecting the access burst, the network (BSC) returns an "immediate assignment" message which includes the following information:

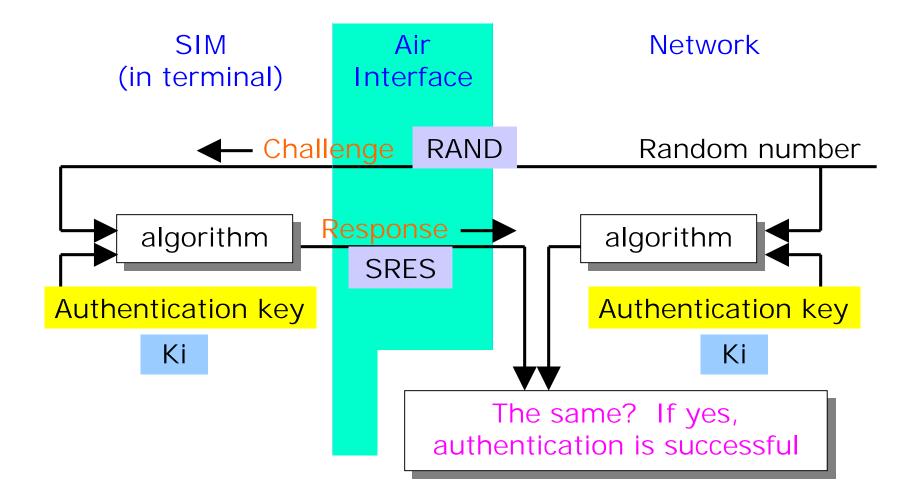
- allocated physical channel (frequency, time slot) in which the assigned signalling channel is located
- timing advance (for correct time slot alignment)

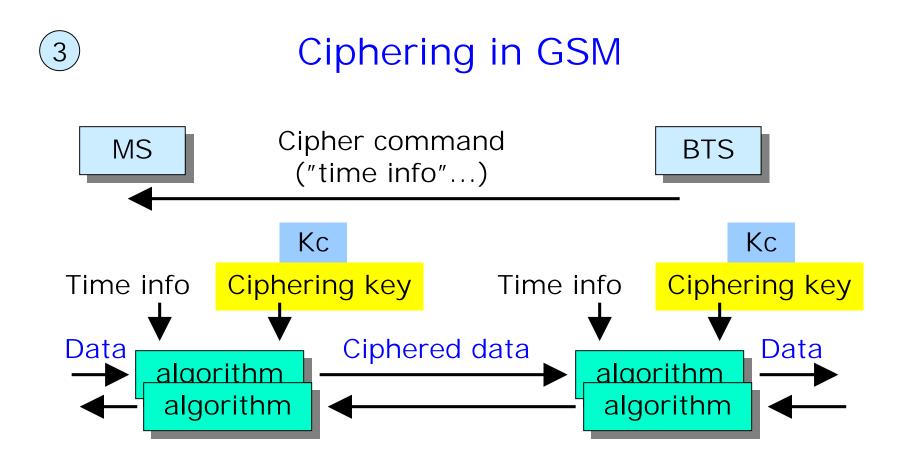
3. The MS now sends a message on the dedicated signalling channel assigned by the network, indicating the reason for performing random access.

Four security measures in GSM

- 1) PIN code (authentication of user using terminal => local security measure, network is not involved)
- 2) SIM authentication (performed by network)
- 3) Ciphering of information sent over air interface
- 4) Usage of TMSI (instead of IMSI) over air interface
 - IMSI = International Mobile Subscriber Identity (globally unique identity)
 - TMSI = Temporary Mobile Subscriber Identity (local and temporary identity)

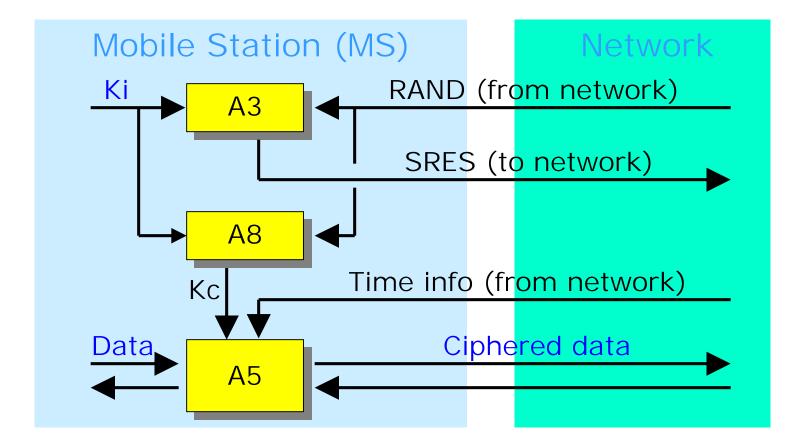
2 Basic principle of user authentication



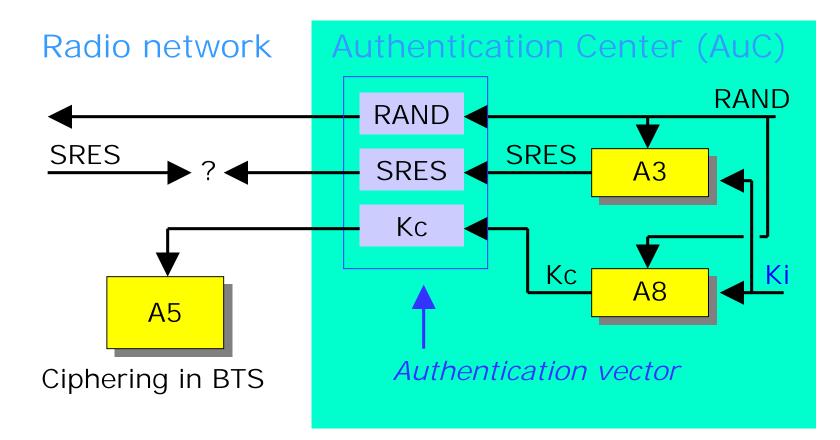


For each call, a new ciphering key (Kc) is generated during authentication both in MS and MSC (in same way as authentication "response").

2 Three security algorithms in GSM (in UMTS many more ...)



2 Three security algorithms in GSM at the network side ...



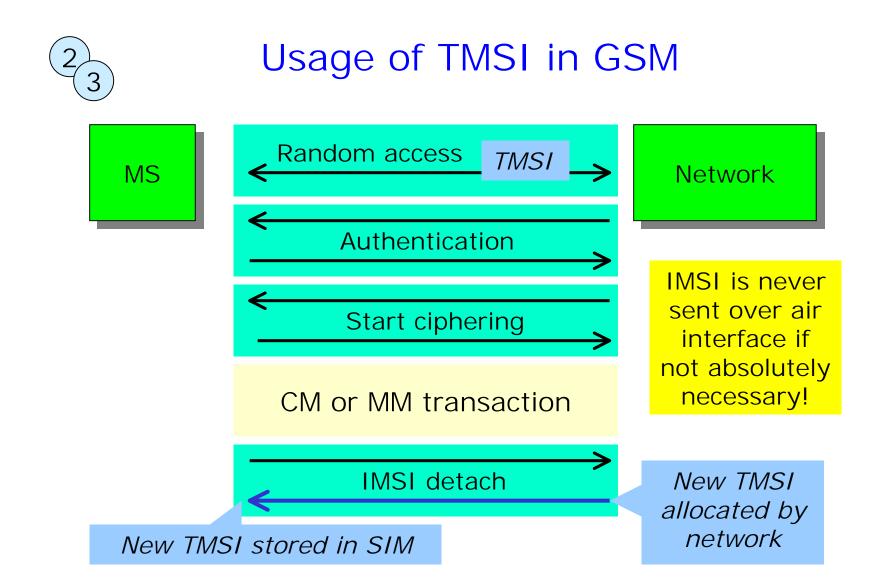


Using output and one or more inputs, it is in practice not possible to calculate "backwards" other input(s) "brute force approach", "extensive search"

Key length in bits (N) is important (in case of brute force approach 2^N calculation attempts may be needed)

Strength of algorithm is that it is secret => bad idea! "security through obscurity"

Better: open algorithm can be tested by engineering community (security through strong algorithm)



4 Connectivity states in GSM/GPRS

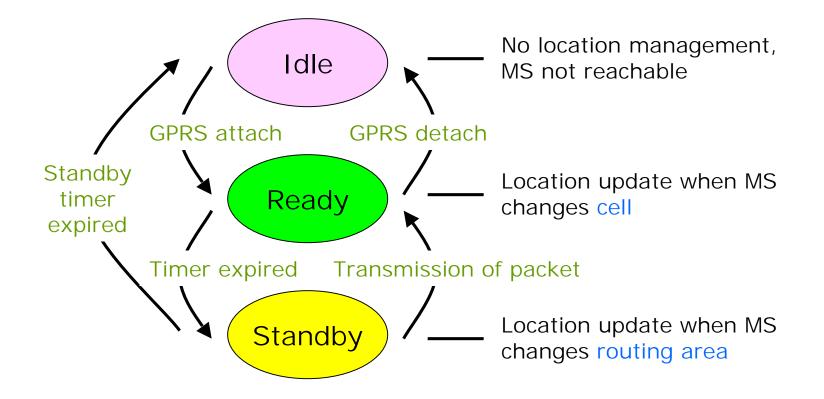
GSM

Disconnected Idle Connected MS is switched off (circuit mode) location updates on LA basis handovers, not location updates

GPRS

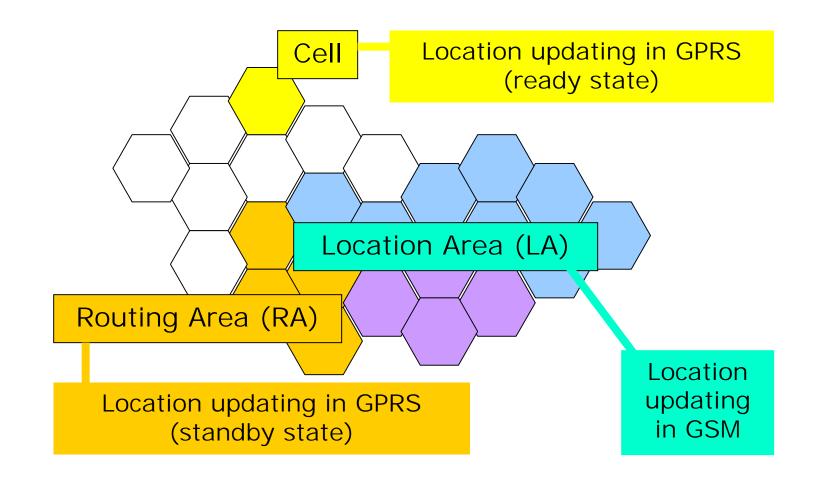
Idle Standby Ready MS is switched off (packet mode) location updates on RA basis location updates on cell basis

4 GPRS connectivity state model



4

MM "areas" in GSM/GPRS



4 Trade-off when choosing LA/RA size

If LA/RA size is very large (e.g. whole mobile network)

- + location updates not needed very often
- paging load is very heavy ,

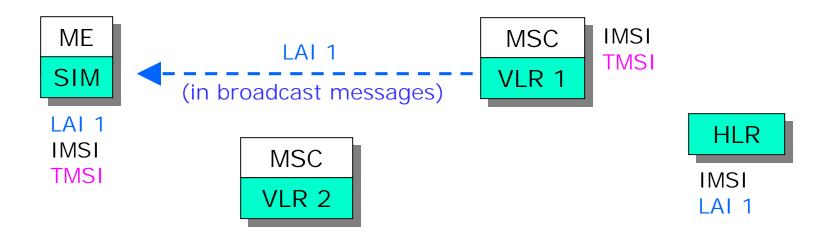
If LA/RA size is very small (e.g. single cell)

- + small paging load
- location updates must be done very often

Affects signalling load

Affects capacity

Case study: GSM location update (1) (most generic scenario)

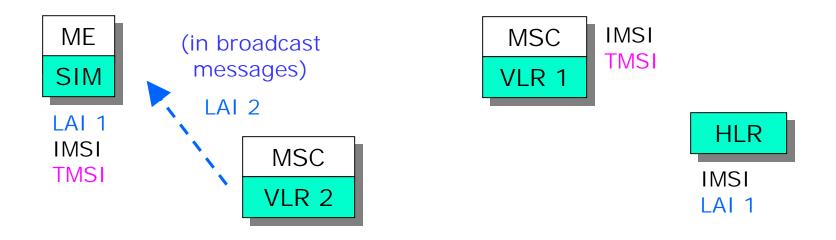


Most recently allocated TMSI and last visited LAI (Location Area ID) are stored in SIM even after switch-off.

After switch-on, MS monitors LAI. If stored and monitored LAI values are the same, no location updating is needed.



GSM location update (2)



Different LAI values => location update required !



SIM sends old LAI and TMSI to VLR 2.

MSC

VLR 2

TMSI

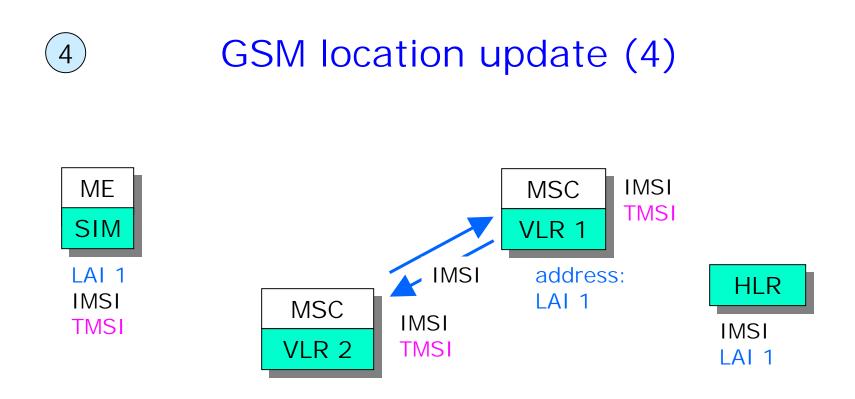
VLR 2 does not recognize TMSI since there is no TMSI-IMSI context. Who is this user?

context

No TMSL - IMSL

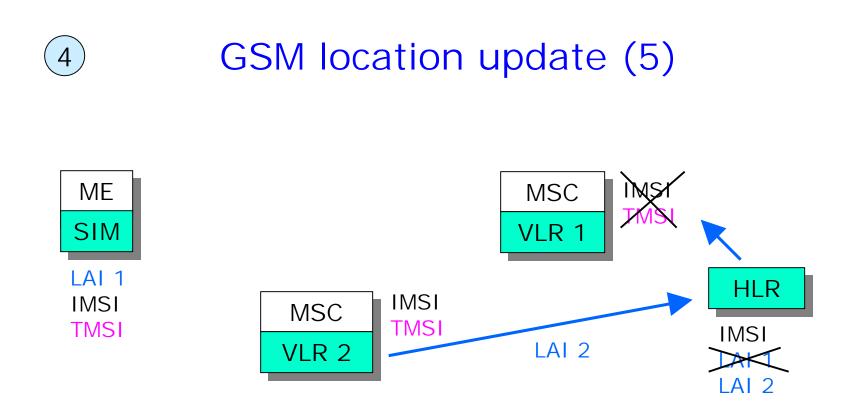
IMSI

LAI 1



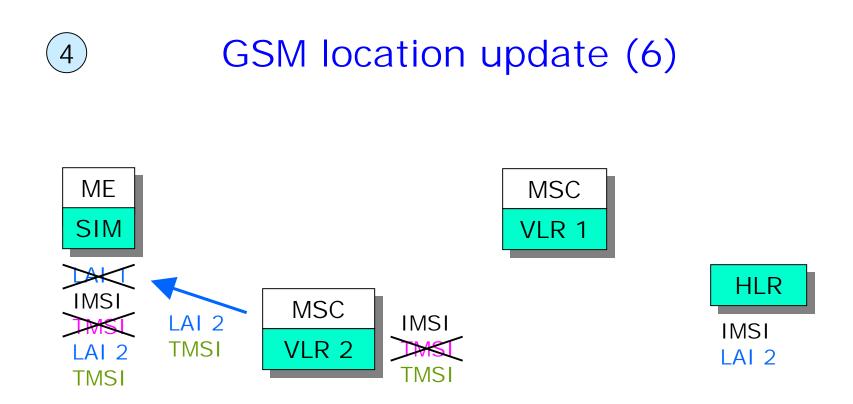
However, VLR 2 can contact VLR 1 (address: LAI 1) and request IMSI.

IMSI is sent to VLR 2.



Important: HLR must be updated (new LAI). If this is not done, incoming calls can not be routed to new MSC/VLR.

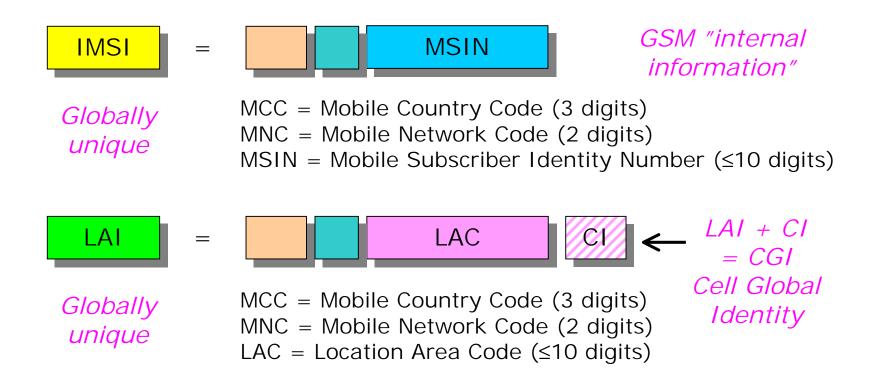
HLR also requests VLR 1 to remove old user data.



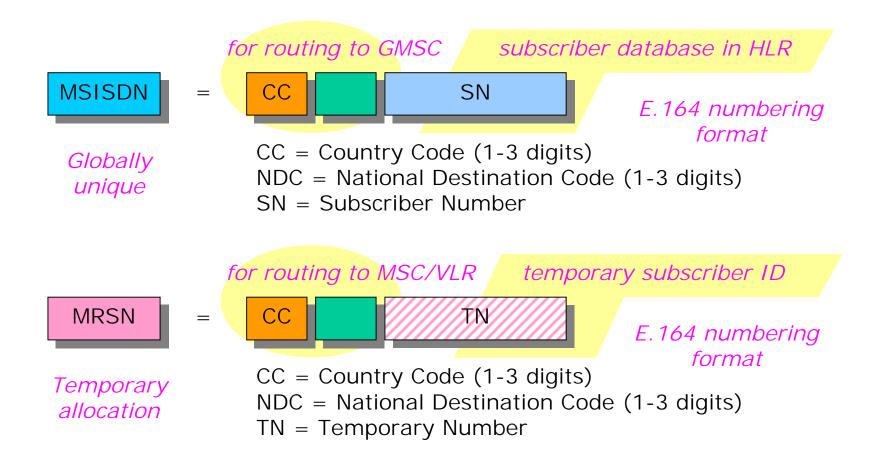
VLR 2 generates new TMSI and sends this to user. User stores new LAI and TMSI safely in SIM.

Location update successful !

GSM identifiers (1)

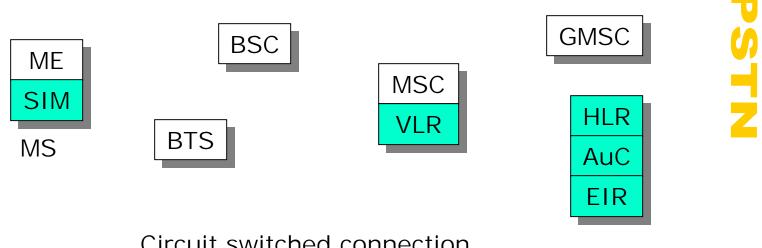


GSM identifiers (2)



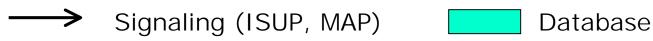


Case study: GSM MTC (1) MTC = mobile terminated call

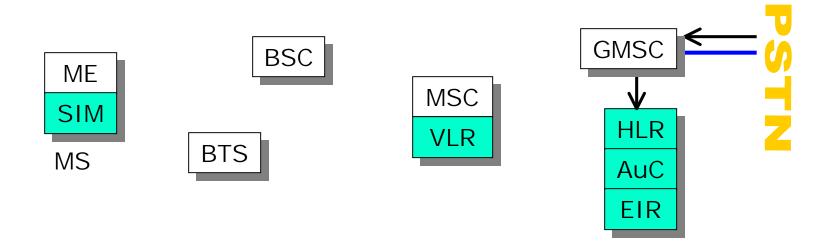


Circuit switched connection

(64 kb/s PCM, 16 kb/s between TRAU and BTS, 13 kb/s encoded speech over air interface)





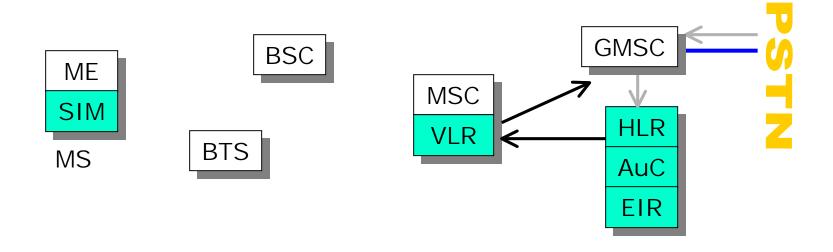


Call is routed to GMSC using MSISDN number of called user (e.g. 040 1234567).

MSISDN number in fact points to database in HLR.

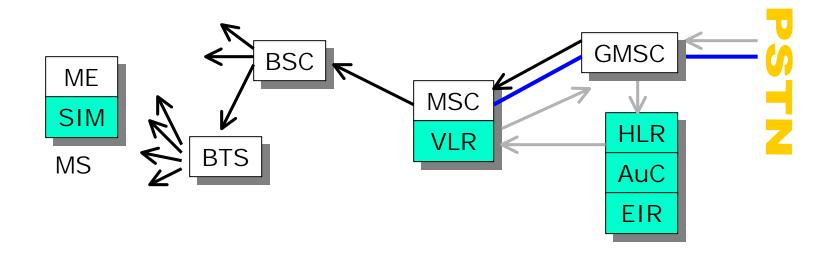
HLR is contacted. Under which MSC/VLR is user?





HLR knows location of Serving MSC/VLR (when user moves to another VLR, this is always recorded in HLR). HLR requests MSRN (roaming number) from VLR. MSRN is forwarded to GMSC.

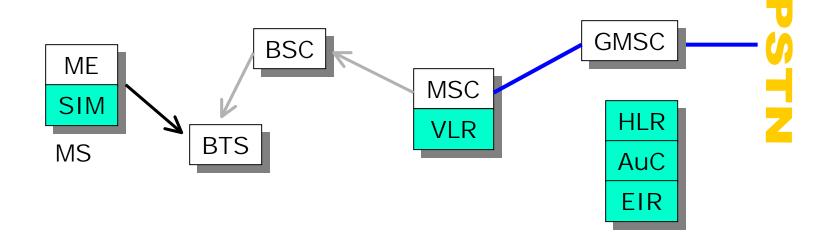




Call can now be routed to Serving MSC/VLR using ISUP (may involve several intermediate switching centers).

MSC/VLR starts paging within Location Area (LA) in which user is located, using TMSI for identification.

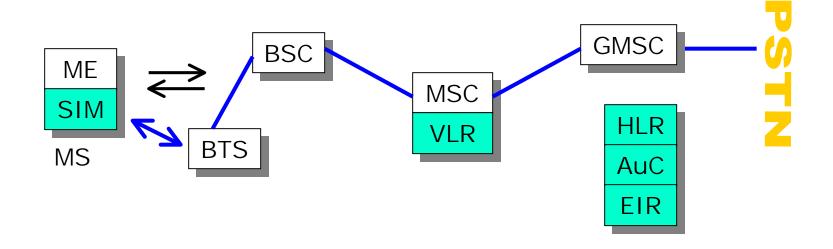




Only the mobile user with the corresponding TMSI responds to the paging.

Using random access procedure, user requests a channel, e.g. SDCCH, for call control signaling.





Signaling channel is set up. After authentication and ciphering procedures, call control signaling continues.

Finally, a GSM traffic channel is set up over the radio interface. The circuit switched connection is now ready.



GPRS attach

Separate or combined GSM/GPRS attach MS registers with an SGSN (authentication...) Location updates possible

PDP context is created

MS is assigned PDP (IP) address < Packet transmission can take place

GPRS detach

PDP context terminated Allocated IP address released In case of dynamic address allocation

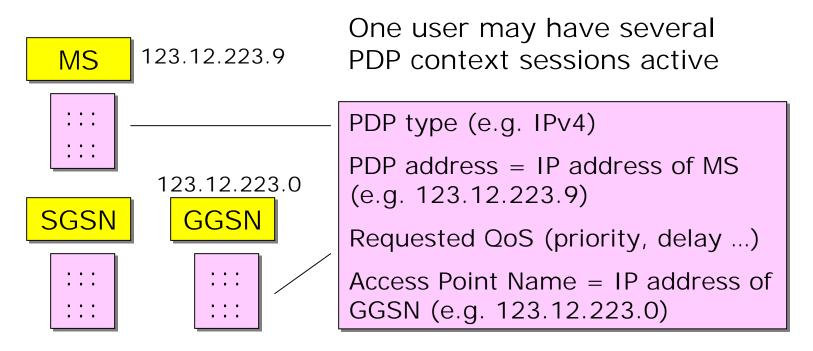
DHCP (Dynamic Host Configuration Protocol)

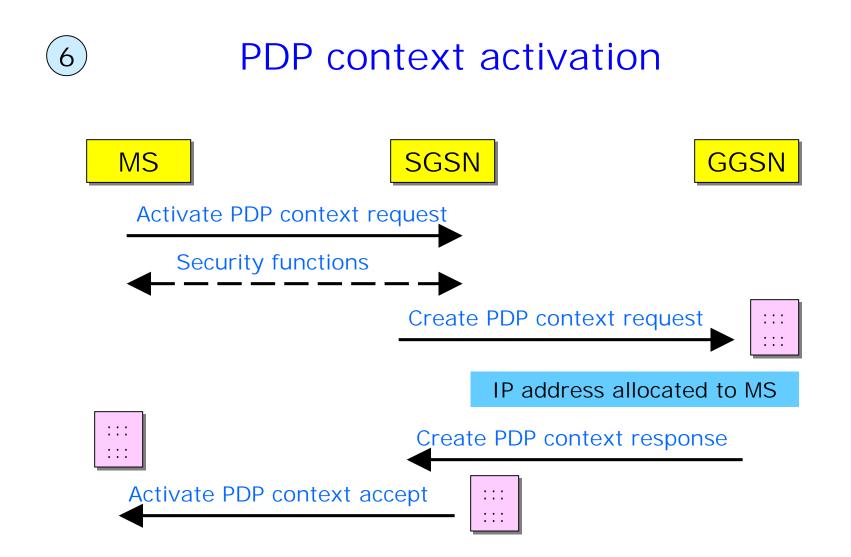
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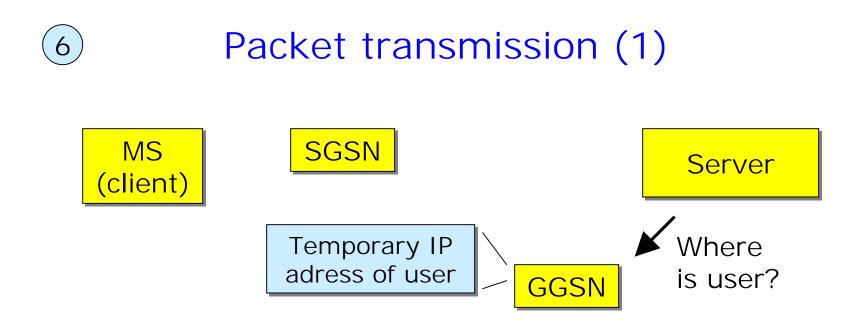
PDP context

PDP context describes characteristics of GPRS session (session = "always on" connection)

PDP context information is stored in MS, SGSN and GGSN



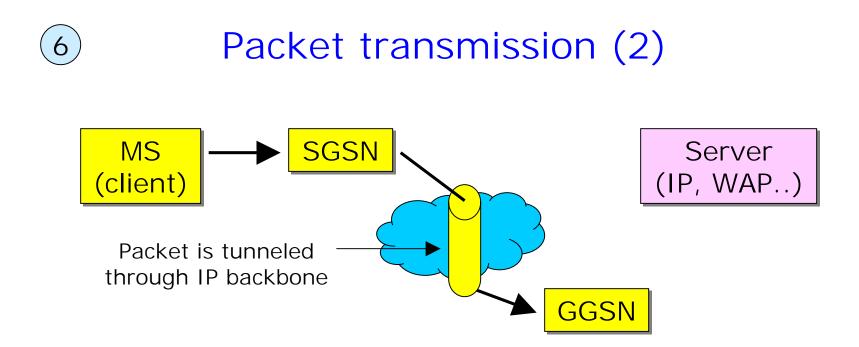




Dynamic IP address allocation has one problem: it is difficult to handle a mobile terminated transaction (external source does not know IP address of MS)

Fortunately, packet services are usually of client-server type

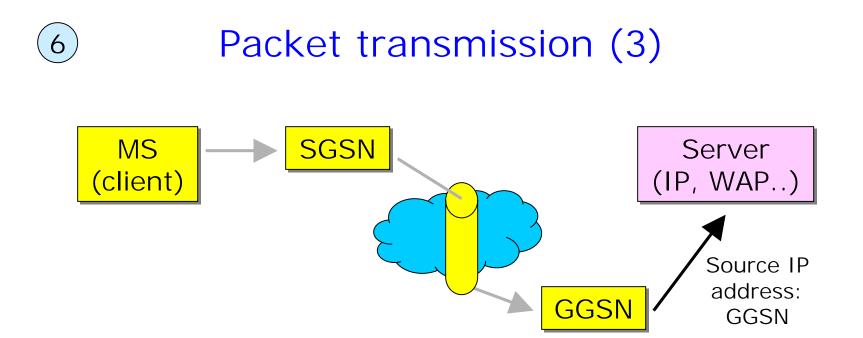
=> MS initiates packet transmission



Packet is sent to SGSN. SGSN sends packet to GGSN through GTP (GPRS Tunneling Protocol) tunnel.

Tunneling = encapsulation of IP packet in GTP packet

IP address	IP address	IP payload	
= APN of GGSN, used for routing through tunnel			

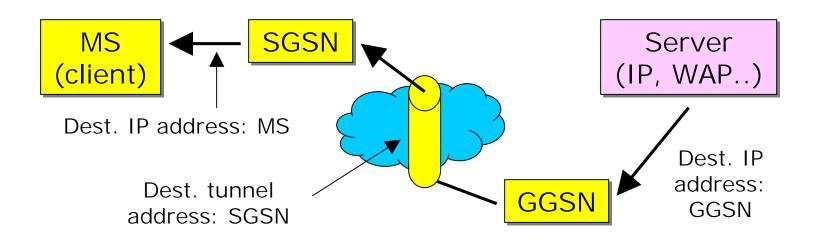


GGSN sends packet through external IP network (i.e. Internet) to the server.

Source IP addr.	Dest. IP addr.	IP payload
GGSN	Server	

Packet transmission (4)

6



Server sends return packet via GGSN, GTP tunnel and SGSN to MS.

Packets from server to MS are always routed via GGSN (since this node has PDP context information).

Further information on GSM/GPRS

Books:

Many good books available (GSM) GPRS is more problematic ...

Web link (GPRS basics):

www.comsoc.org/livepubs/surveys/public/3q99issue/ bettstetter.html